

# NAG Fortran Library Routine Document

## F01BVF

**Note:** before using this routine, please read the Users' Note for your implementation to check the interpretation of *bold italicised* terms and other implementation-dependent details.

### 1 Purpose

F01BVF transforms the generalized symmetric-definite eigenproblem  $Ax = \lambda Bx$  to the equivalent standard eigenproblem  $Cy = \lambda y$ , where  $A$ ,  $B$  and  $C$  are symmetric band matrices and  $B$  is positive-definite.  $B$  must have been decomposed by F01BUF.

### 2 Specification

```
SUBROUTINE F01BVF(N, MA1, MB1, M3, K, A, IA, B, IB, V, IV, W, IFAIL)
INTEGER          N, MA1, MB1, M3, K, IA, IB, IV, IFAIL
real           A(IA,N), B(IB,N), V(IV,M3), W(M3)
```

### 3 Description

$A$  is a symmetric band matrix of order  $n$  and bandwidth  $2m_A + 1$ . The positive-definite symmetric band matrix  $B$ , of order  $n$  and bandwidth  $2m_B + 1$ , must have been previously decomposed by F01BUF as  $ULDL^T U^T$ . F01BVF applies  $U$ ,  $L$  and  $D$  to  $A$ ,  $m_A$  rows at a time, restoring the band form of  $A$  at each stage by plane rotations. The parameter  $k$  defines the change-over point in the decomposition of  $B$  as used by F01BUF and is also used as a change-over point in the transformations applied by this routine. For maximum efficiency,  $k$  should be chosen to be the multiple of  $m_A$  nearest to  $n/2$ . The resulting symmetric band matrix  $C$  is overwritten on  $A$ . The eigenvalues of  $C$ , and thus of the original problem, may be found using F08HEF (SSBTRD/DSBTRD) and F08JFF (SSTERF/DSTERF). For selected eigenvalues, use F08HEF (SSBTRD/DSBTRD) and F08JF (SSTEBZ/DSTEBZ).

### 4 References

Crawford C R (1973) Reduction of a band-symmetric generalized eigenvalue problem *Comm. ACM* **16** 41–44

### 5 Parameters

- 1: N – INTEGER *Input*  
*On entry:*  $n$ , the order of the matrices  $A$ ,  $B$  and  $C$ .
- 2: MA1 – INTEGER *Input*  
*On entry:*  $m_A + 1$ , where  $m_A$  is the number of non-zero super-diagonals in  $A$ . Normally  $MA1 \ll N$ .
- 3: MB1 – INTEGER *Input*  
*On entry:*  $m_B + 1$ , where  $m_B$  is the number of non-zero super-diagonals in  $B$ .  
*Constraint:*  $MB1 \leq MA1$ .
- 4: M3 – INTEGER *Input*  
*On entry:* the value of  $3m_A + m_B$ .

- 5: K – INTEGER *Input*
- On entry:*  $k$ , the change-over point in the transformations. It must be the same as the value used by F01BUF in the decomposition of  $B$ .
- Suggested value:* the optimum value is the multiple of  $m_A$  nearest to  $n/2$ .
- Constraint:*  $MB1 - 1 \leq K \leq N$ .
- 6: A(IA,N) – *real* array *Input/Output*
- On entry:* the upper triangle of the  $n$  by  $n$  symmetric band matrix  $A$ , with the diagonal of the matrix stored in the  $(m_A + 1)$ th row of the array, and the  $m_A$  super-diagonals within the band stored in the first  $m_A$  rows of the array. Each column of the matrix is stored in the corresponding column of the array. For example, if  $n = 6$  and  $m_A = 2$ , the storage scheme is
- |          |          |          |          |          |          |
|----------|----------|----------|----------|----------|----------|
| *        | *        | $a_{13}$ | $a_{24}$ | $a_{35}$ | $a_{46}$ |
| *        | $a_{12}$ | $a_{23}$ | $a_{34}$ | $a_{45}$ | $a_{56}$ |
| $a_{11}$ | $a_{22}$ | $a_{33}$ | $a_{44}$ | $a_{55}$ | $a_{66}$ |
- Elements in the top left corner of the array need not be set. The following code assigns the matrix elements within the band to the correct elements of the array:
- ```

      DO 20 J = 1, N
        DO 10 I = MAX(1,J-MA1+1), J
          A(I-J+MA1,J) = matrix(I,J)
        10 CONTINUE
      20 CONTINUE

```
- On exit:*  $A$  is overwritten by the corresponding elements of  $C$ .
- 7: IA – INTEGER *Input*
- On entry:* the first dimension of the array  $A$  as declared in the (sub)program from which F01BVF is called.
- Constraint:*  $IA \geq MA1$ .
- 8: B(IB,N) – *real* array *Input/Output*
- On entry:* the elements of the decomposition of matrix  $B$  as returned by F01BUF.
- On exit:* the elements of  $B$  will have been permuted.
- 9: IB – INTEGER *Input*
- On entry:* the first dimension of the array  $B$  as declared in the (sub)program from which F01BVF is called.
- Constraint:*  $IB \geq MB1$ .
- 10: V(IV,M3) – *real* array *Workspace*
- 11: IV – INTEGER *Input*
- On entry:* the first dimension of the array  $V$  as declared in the (sub)program from which F01BVF is called.
- Constraint:*  $IV \geq m_A + m_B$ .
- 12: W(M3) – *real* array *Workspace*
- 13: IFAIL – INTEGER *Input/Output*
- On entry:* IFAIL must be set to 0, -1 or 1. Users who are unfamiliar with this parameter should refer to Chapter P01 for details.
- On exit:* IFAIL = 0 unless the routine detects an error (see Section 6).

For environments where it might be inappropriate to halt program execution when an error is detected, the value  $-1$  or  $1$  is recommended. If the output of error messages is undesirable, then the value  $1$  is recommended. Otherwise, for users not familiar with this parameter the recommended value is  $0$ . **When the value  $-1$  or  $1$  is used it is essential to test the value of IFAIL on exit.**

## 6 Error Indicators and Warnings

If on entry IFAIL =  $0$  or  $-1$ , explanatory error messages are output on the current error message unit (as defined by X04AAF).

Errors or warnings detected by the routine:

IFAIL =  $1$

On entry, MB1 > MA1.

## 7 Accuracy

In general the computed system is exactly congruent to a problem  $(A + E)x = \lambda(B + F)x$ , where  $\|E\|$  and  $\|F\|$  are of the order of  $\epsilon\kappa(B)\|A\|$  and  $\epsilon\kappa(B)\|B\|$  respectively, where  $\kappa(B)$  is the condition number of  $B$  with respect to inversion and  $\epsilon$  is the *machine precision*. This means that when  $B$  is positive-definite but not well-conditioned with respect to inversion, the method, which effectively involves the inversion of  $B$ , may lead to a severe loss of accuracy in well-conditioned eigenvalues.

## 8 Further Comments

The time taken by the routine is approximately proportional to  $n^2m_B^2$  and the distance of  $k$  from  $n/2$ , e.g.,  $k = n/4$  and  $k = 3n/4$  take 502% longer.

When  $B$  is positive-definite and well-conditioned with respect to inversion, the generalized symmetric eigenproblem can be reduced to the standard symmetric problem  $Py = \lambda y$  where  $P = L^{-1}AL^{-T}$  and  $B = LL^T$ , the Cholesky factorization.

When  $A$  and  $B$  are of band form, especially if the bandwidth is small compared with the order of the matrices, storage considerations may rule out the possibility of working with  $P$  since it will be a full matrix in general. However, for any factorization of the form  $B = SS^T$ , the generalized symmetric problem reduces to the standard form

$$S^{-1}AS^{-T}(S^T x) = \lambda(S^T x)$$

and there does exist a factorization such that  $S^{-1}AS^{-T}$  is still of band form (see Crawford (1973)). Writing

$$C = S^{-1}AS^{-T} \quad \text{and} \quad y = S^T x$$

the standard form is  $Cy = \lambda y$  and the bandwidth of  $C$  is the maximum bandwidth of  $A$  and  $B$ .

Each stage in the transformation consists of two phases. The first reduces a leading principal sub-matrix of  $B$  to the identity matrix and this introduces non-zero elements outside the band of  $A$ . In the second, further transformations are applied which leave the reduced part of  $B$  unaltered and drive the extra elements upwards and off the top left corner of  $A$ . Alternatively,  $B$  may be reduced to the identity matrix starting at the bottom right-hand corner and the extra elements introduced in  $A$  can be driven downwards.

The advantage of the  $ULDL^T U^T$  decomposition of  $B$  is that no extra elements have to be pushed over the whole length of  $A$ . If  $k$  is taken as approximately  $n/2$ , the shifting is limited to halfway. At each stage the size of the triangular bumps produced in  $A$  depends on the number of rows and columns of  $B$  which are eliminated in the first phase and on the bandwidth of  $B$ . The number of rows and columns over which these triangles are moved at each step in the second phase is equal to the bandwidth of  $A$ .

In this routine,  $A$  is defined as being at least as wide as  $B$  and must be filled out with zeros if necessary as it is overwritten with  $C$ . The number of rows and columns of  $B$  which are effectively eliminated at each stage is  $m_A$ .



```

        WRITE (NOUT,99999) 'ssbtrd', INFO
    ELSE
*
        ABSTOL = 0.0e0
        READ (NIN,*) M1, M2
*
        CALL sstebz('I','E',N,0.0e0,0.0e0,M1,M2,ABSTOL,D,E,M,NSPLIT,
+           R,IBLOCK,ISPLIT,WORK,IWORK,INFO)
        IF (INFO.NE.0) THEN
            WRITE (NOUT,99999) 'sstebz', INFO
        ELSE
            WRITE (NOUT,*)
            WRITE (NOUT,*) 'Selected eigenvalues'
            WRITE (NOUT,99998) (R(I),I=1,M)
        END IF
    END IF
END IF
STOP
*
99999 FORMAT (1X,'INFO from ',A6,' = ',I3)
99998 FORMAT (1X,8F9.4)
END

```

## 9.2 Program Data

F01BVF Example Program Data

```

9  2  2
11
12  12
13  13
14  14
15  15
16  16
17  17
18  18
19  19
101
22  102
23  103
24  104
25  105
26  106
27  107
28  108
29  109
1  3

```

### 9.3 Program Results

F01BVF Example Program Results

Selected eigenvalues  
-0.2643 -0.1530 -0.0418

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